

# Projection

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# Model

$$\begin{aligned}c_t^{-\nu} &= \mathbb{E}_t \left[ \beta c_{t+1}^{-\nu} \alpha z_{t+1} k_{t+1}^{\alpha-1} \right] \\c_t + k_{t+1} &= z_t k_t^\alpha \\ \ln(z_{t+1}) &= \rho \ln(z_t) + \varepsilon_{t+1} \\ \varepsilon_{t+1} &\sim N(0, \sigma^2) \\ k_1, z_1 &\text{ given}\end{aligned}$$

$k_t$  is beginning-of-period  $t$  capital stock

# Projection Methods

**True rational expectations solution:**

$$c_t = c(k_t, z_t)$$

- Why a difficult problem to find  $c(k_t, z_t)$ ?

## Approximation:

- structural parameters  $(\alpha, \beta, \rho, \sigma)$  have fixed numerical values (thus not included as arguments in function)

$$c_t = c(k_t, z_t) \approx P_n(k_t, z_t; \eta_n)$$

- $P_n(\cdot)$ : from class of approximating functions
  - such as polynomials or splines
- $n$  is fixed  $\implies$  solve for  $\eta_n$ , a *finite-dimensional* object

# Which equations to use?

- goal: solve for  $P_n(k_t, z_t; \eta_n) \approx c(k_t, z_t)$ ,
  - i.e.,  $N_n$  elements of  $\eta_n$
  - $k(k_t, z_t)$  implicitly defined by budget constraint
- One first-order equation left, namely Euler equation
  - this is a different equation at each point in the state space
  - $\implies$  plenty of equations

# Which equations to use?

- Equations at  $M$  grid points  $\{k_i, z_i\}$  with  $M \geq N_n$

$$P_n(k_i, z_i; \eta_n)^{-\nu} =$$

$$E \begin{bmatrix} \alpha \beta \times \\ P_n(\{\mathbf{k}'\}, \{\mathbf{z}'\}; \eta_n)^{-\nu} \times \\ \{\mathbf{z}'\} \times \\ (\{\mathbf{k}'\})^{\alpha-1} \end{bmatrix}$$

# Which equations to use?

- **Goal:**  $\forall$  grid point get 1 equation with  $\eta_n$  as only unknown

$$E \begin{bmatrix} P_n(k_i, z_i; \eta_n)^{-\nu} = \\ \alpha\beta \times \\ P_n(\mathbf{k}', \mathbf{z}'; \eta_n)^{-\nu} \times \\ \mathbf{z}' \times \\ (\mathbf{k}')^{\alpha-1} \end{bmatrix}$$

- Note that  $k_i$  and  $z_i$  are known

# Which equations to use?

$$E \left[ \begin{array}{c} P_n(k_i, z_i; \eta_n)^{-\nu} = \\ \alpha \beta \times \\ P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \varepsilon'\}; \eta_n)^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \varepsilon'\} \times \\ (z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1} \end{array} \right]$$

# How to deal with expectations operator?

Let  $\{\omega_j, \zeta_j\}_{j=1}^J$  be the Hermite Gaussian quadrature nodes

$$P_n(k_i, z_i; \eta_n)^{-\nu} =$$

$$\sum_{j=1}^J \left[ \begin{array}{c} \alpha\beta \times \\ P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n)^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

# Define error terms

$$e(k_i, z_i; \eta_n) = P_n(k_i, z_i; \eta_n)^{-\nu} -$$

$$\sum_{j=1}^J \left[ \begin{array}{c} \alpha \beta \times \\ P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n)^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (z_i k_i^\alpha - P_n(k_i, z_i; \eta_n))^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

# How to find coefficients of approximation?

- True rational expect. solution gives zero error term  $\forall (k_i, z_i)$
- Thus, choose  $\eta_n$  such that error terms are as small as possible.
- **Collacation** ( $M = N_n$ ): Use equation solver to get errors exactly equal to zero on grid
- **Galerkin** ( $M > N_n$ ): Use minimization routine (and possibly smart weighting of error terms)

# Different types of approximating functions

- $P_n(k_i, z_i; \eta_n)$  could be polynomial or spline
- dimension  $\eta_n$  usually higher for splines
  - may make eq. solver/minimization less appropriate
  - use iteration scheme instead

# How to find coefficients of approximation?

- ① Equation solver or minimization routine
- ② Iteration procedures
  - ① fixed-point iteration
  - ② time iteration

# Iterating versus eq. solver/minimization

- Advantage:
  - less of a black box
  - can deal with many coefficients
    - e.g. when spline is used
  - some iteration schemes are guaranteed to converge
    - under some regularity conditions
- Disadvantage:
  - does not use information on how best to update

# Iteration procedure: Construct Grid

- Construct a grid with nodes for  $k$  and  $z$
- At the nodes construct the basis functions of  $P_n(k, z; \eta_n)$ .
- For example, if

$$P_n(k, z; \eta_n) = \eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2$$

then construct the matrix (where subscripts denote grid numbers)

$$X = \begin{bmatrix} 1 & k_1 & z_1 & k_1^2 & k_1 z_1 & z_1^2 \\ 1 & k_2 & z_2 & k_2^2 & k_2 z_2 & z_2^2 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & k_M & z_M & k_M^2 & k_M z_M & z_M^2 \end{bmatrix}$$

and calculate  $(X'X)^{-1} X'$

# Iteration procedure: Construct Grid

- **Chebyshev nodes:** Using Chebyshev nodes is important. This ensures uniform convergence. With equidistant nodes it is possible that the oscillations between grid point explode as the order of the polynomial increases.
- **Chebyshev polynomials:** If you have (i) no problems finding initial conditions and (ii) only low-order approximations so that calculating the inverse of  $X'X$  can be done accurately, then you can use regular polynomials. Orthogonal Chebyshev polynomials can overcome these problems. They ensure that  $X'X$  is diagonal (and trivial to invert). This does require scaling of the state variables so they are between  $-1$  and  $1$ .

# Fixed-point Iteration

The value of  $\eta_n$  used in the  $q^{\text{th}}$  iteration is referred to as  $\eta_n^q$ . Follow the following iteration scheme until convergence

- At each grid point:
  - Calculate the RHS of the Euler equation using the latest value for  $\eta_n$ , i.e.,  $\eta_n^{q-1}$
  - Use RHS to calculate  $c_i$ , value for  $c$  at  $i^{\text{th}}$  grid point
- Use values for  $c_i$  to obtain an estimate for  $\eta_n$ ,  $\hat{\eta}_n^q$ 
  - Polynomial: run a regression to get  $\hat{\eta}_n^q$
  - Spline: the values of  $c$  at the nodes are the new values of  $\eta_n$
- Let  $\eta_n^q = \lambda \hat{\eta}_n^q + (1 - \lambda) \eta_n^{q-1}$

# Fixed-point Iteration

- **Step 1: Calculate current consumption values implied by  $\eta_n^{j-1}$  at each grid point**
  - Use  $\eta_n^{q-1}$  to calculate  $k' = z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1})$
  - Use  $\eta_n^{q-1}$  to calculate  $c' = P_n(k', z'; \eta_n^{q-1})$
  - Then, get  $c_i$  from

$$(c_i)^{-\nu} =$$

$$\sum_{j=1}^J \left[ \begin{array}{c} \alpha\beta \times \\ P_n(z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1}), \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ \left( z_i k_i^\alpha - P_n(k_i, z_i; \eta_n^{q-1}) \right)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

# Fixed-point iteration

**Step 2: Get new estimate for  $\eta_n$  by running a projection step**

- Let  $Y = [c_1, c_2, \dots, c_M]'$
- If

$$P_n(k, z; \eta_n) = \eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2$$

then

$$\hat{\eta}_n^q = (X'X)^{-1} X'Y$$

# Fixed-point iteration

**Step 2: Get new estimate for  $\eta_n$  by running a projection step**

- If

$$P_n(k, z; \eta_n) = \exp \left( \eta_{0,n} + \eta_{k,n}k + \eta_{z,n}z + \eta_{kk}k^2 + \eta_{kz}kz + \eta_{zz}z^2 \right)$$

then

$$\hat{\eta}_n^q = (X'X)^{-1} X' \ln(Y)$$

- no stochastic error term  $\implies$  ok to take  $\ln$  of LHS & RHS

# Fixed-point iteration

## Step 3: Update $\eta_n$

$$\eta_n^q = \lambda \widehat{\eta}_n^q + (1 - \lambda) \eta_n^{q-1} \quad \text{for } 0 < \lambda \leq 1$$

- Fixed-point iteration does not always converge
  - Choosing a lower value of  $\lambda$ :
    - convergence more likely
    - slows down algorithm if lower value not needed for convergence
- Alternative is **time iteration**

# Time Iteration

- At each grid point use  $\eta_n^{q-1}$  only for *next period's* choices
- Again solve for  $c_i$  at each grid point
  - this is now a bit trickier (non-linear problem)
- Get  $n_n^q$  as with fixed-point iteration
  - guaranteed to converge without dampening (under regularity conditions)

# Time Iteration - solving for $c$

Solve  $c_i$  from following non-linear equation

$$(c_i)^{-\nu} = \sum_{j=1}^J \left[ \begin{array}{c} \alpha\beta \times \\ P_n(z_i k_i^\alpha - c_i, \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (z_i k_i^\alpha - c_i)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

# Time Iteration

- Natural interpretation for  $\eta_n^{q-1}$  and  $\eta_n^q$ , namely
  - $\eta_n^{q-1}$  is tomorrow's policy function and
  - $\eta_n^q$  is today's policy function
- Time iteration is reliable and convergent
  - (the proof is related to the convergence of value function iteration, which uses the same idea)

# Fixed-point versus time iteration

- Fixed-point iteration uses  $\eta_n^{q-1}$  for *all* terms on the RHS, i.e., both next period's consumption choice and today's capital choice
- Time iteration uses  $\eta_n^{q-1}$  only to evaluate next period's consumption
- The structure of time iteration mimics the choice of value function iteration:
  - next period's behavior described by previous solution for value function
  - Bellman equation used to solve for choice of  $c$  and  $k$  *simultaneously*

# Endogenous grid points

- Simple idea: construct grid for  $k'$  instead of a grid for  $k$
- Instead of solving for the choice  $k'$  given  $k$ , we now solve for the value of  $k$  that would have led to the choice  $k'$
- In both cases you end up at each grid point with a set of values for  $k$  and a set of corresponding values for  $k'$ .
- Terminology is a bit confusing: the grid itself is exogenous and fixed but it is for an endogenous variable
- You can use endogenous grid points both with fixed-point and with time iteration
- The added value with time iteration lies in getting rid of the non-linear problem of solving for today's choices

# Endogenous grid points and time iteration

- Time iteration  $\implies$ 
  - use  $\eta_n^{q-1}$  for tomorrow's choices and
  - use  $\eta_n^q$  only for today's choices (which show up on both sides of the policy function)
- Then, get  $c_i$  from

# Endogenous grid points and time iteration

$$(c_i)^{-\nu} = \sum_{j=1}^J \left[ \begin{array}{c} \alpha\beta \times \\ P_n(k'_i, \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\}; \eta_n^{q-1})^{-\nu} \times \\ \exp\{\rho \ln(z_i) + \sqrt{2}\sigma\zeta_j\} \times \\ (k'_i)^{\alpha-1} \\ \omega_j / \sqrt{\pi} \end{array} \right]$$

and  $k_i$  from

$$k'_i + c_i = z_i k^\alpha$$

# Perturbation versus projection

- Nondifferentiabilities
  - impossible for perturbation
- Large number of state variables
  - difficult for projection
- Constructing the grid can be difficult
  - apriori hard to know what sensible points are
  - some calculations may not be well defined everywhere

# Perturbation versus projection

- Global versus local
  - Projection designed to be global method
  - Perturbation designed to be local method
    - but could give accurate global approximation
    - question is whether (lower-order) derivatives at perturbation point capture global behavior

# References

- Judd, K. L., 1998, Numerical Methods in Economics, The MIT Press.
- Rendahl, P., 2006, Inequality constraints in recursive economies.
  - shows that time-iteration converges even in the presence of inequality constraints